

DM sizing model and cost plan for construction and operations. **William O'Mullane, Richard Dubois, Michelle Butler, Kian Tat Lim** 2021-03-19

# **1 Introduction**

This document presents the simplified sizing model for Rubin Observatory data management in Section [6.1](#page-8-0) based on detailed sizing presented in Section [7](#page-10-0). Section [2](#page-0-0) presents a very high level budget summary for DM hardware which was used for [LCR-2148.](https://project.lsst.org/groups/ccb/node/3889) More interesting now is the build up at the USDF in pre-operations which is shown in Section [3](#page-1-0) and the full operations estimates in Section [4](#page-2-0)

This version is in agreement with SLAC on the parameters for CPU and disk price fall as well as CPU cost etc.

# <span id="page-0-0"></span>**2 Construction budget**

<span id="page-0-1"></span>A high level bottom line is given in Table [1.](#page-0-1) The remainder of the document is all the details that went into that.



Table 1: This table pulls together all the information in a high level summary - in this table Xeon pricing is used since that is the more expensive but better known option. Price factors, defined in Table [27](#page-8-1) are applied post 2020.

We have applied a modest cost reduction assuming that processors and disks get a little cheaper - that percentage is given in Table [27](#page-8-1) along with many other parameters. Table [27](#page-8-1) also contains the number of nodes we assume to need for Qserv.



Specific costs for storage are detailed in Table [28](#page-9-0) and for compute in Table [29](#page-9-1) the following budgets can be considered. The detailed annual purchasing based on those prices is given for storage in Table [11](#page-4-0) and for compute in Table [10](#page-4-1).

# <span id="page-1-0"></span>**3 Pre-Operations budget estimate**

<span id="page-1-1"></span>In this section we estimate the ramp up of the USDF to be ready for start of operations. This means having compute for commissioning data as well as developer services and a small science platform in place. It is very similar to the construction needs. The summary is in Table [2.](#page-1-1)





Currently we assume exactly the construction profile for storage plus space for datasets currently held at NCSA. This is presented in Table [4.](#page-1-2)

<span id="page-1-3"></span>The compute is a little different and uses Rome which is captured in Table [3](#page-1-3).

Table 3: Preoperations compute build up at USDF, cores we need to purchase per year.



Table 4: Preoperations storage to be purchaed each year

<span id="page-1-2"></span>





# <span id="page-2-0"></span>**4 Operations budget estimate**

Based on the needs in Table [25](#page-7-0) and the costs in Table [18](#page-5-0) and Table [29](#page-9-1) we get the estimate presented in Table [5](#page-2-1). In Table [5](#page-2-1) we should note that IN2P3 do 50% and UKDF do 25% of the processing so we reduce the processing cost by three quarters. This does not reduce the storage cost.

Table 5: This table adds the USDF (see Table [7](#page-3-0)) and Table [8](#page-3-1) to give total operations hardware costs.

<span id="page-2-1"></span>

Again in Table [5](#page-2-1) we assume IN2P3 do 50% of processing (see Table [7\)](#page-3-0). We have applied a compounded modest cost reduction assuming that processors and disks get a little cheaper - that percentage is given in Table [27.](#page-8-1)

It must be noted that the price of disk and tape have a profound effect over 10 years. We have been fairly conservative on the base prices in Table [28.](#page-9-0) An even bigger effect is in the compounding of the presumed fall in storage cost. Here we have used an extremely conservative 5% per year (Table [27\)](#page-8-1) - changing this to 15% halves the cumulative ops estimate, setting it to 10% brings the total down by about 30%.

# **4.1 Cloud costs**

In addition there are some cloud costs. We run certain jobs and host websites on Amazon and Google. In operations the validation team may also wish to run simulations on cloud resources. This estimate is in Table [6](#page-2-2).

> Table 6: We have on going cloud costs and assume some other activities may be on cloud in the future we make an estimate of those costs here.

<span id="page-2-2"></span>





More details on the inputs are in Section [5.1](#page-4-2)

## **4.2 US and Chile**

While Table [5](#page-2-1) present the total ops cost for Rubin Observatory a fraction of this is in Chile and would potentially remain an NSF cost in operations. Table [7](#page-3-0) presents just the US Data Facility budget and Table [8](#page-3-1) presents the Chile budget.

> Table 7: This table pulls together all the information in a high level summary for USDF operations - in this table Rome pricing(see Table [14](#page-5-1)). Price factors, defined in Table [27](#page-8-1) are applied in all cases - other input values come from Table [25,](#page-7-0) Table [19](#page-5-2).

<span id="page-3-0"></span>

<span id="page-3-1"></span>Note that in Table [8](#page-3-1) a fraction of the USDF storage is considered, big enough for Raws and a cache of products. It also does not include a Qserv. See also Table [26](#page-8-2)

> Table 8: "This table pulls together all the information in a high level summary for Chile operations - in this table Rome pricing(see Table [17\)](#page-5-3) is used Price factors, defined in Table [27](#page-8-1) are applied in all cases - other input values come from Table [25](#page-7-0), Table [20.](#page-5-4)



# **5 Cost details**

The summary table (Table [1](#page-0-1)) uses Xeon pricing for compute as shown in Table [9](#page-4-3).







<span id="page-4-3"></span><span id="page-4-1"></span>An alternative architecture would be Rome - SLAC have chosen this for the Ops pricing, Ta-ble [10](#page-4-1) gives the price of compute based on Rome -small and large. Rome large are used in the operations calculations.

Table 10: Implementation with AMD Rome (we have no good proce for these reallly)



<span id="page-4-0"></span>Table [11](#page-4-0) gives the price of storage using all types that we need. This would be needed regardless of the compute chosen.

Table 11: Total storage cost estimate



<span id="page-4-4"></span>Table [12](#page-4-4) gives the annual cost of hosting compute at NCSA for construction. This includes purchasing racks to house new nodes etc.

> Table 12: Overheads(NCSA) per year based on number of cores in Table [27](#page-8-1) and costs in Table [30](#page-10-1) assuming Xeon density from Table [29](#page-9-1).



## <span id="page-4-2"></span>**5.1 Ops Cost details**

Table [13](#page-5-5) gives the price of compute based on Xeons. This is broken down further for US in Table [15](#page-5-6) and Chile in Table [16](#page-5-7). However the Ops costing for SLAC was done using Table [14.](#page-5-1)



Table [18](#page-5-0) gives the price of storage using all types that we need. This is broken down further for US in Table [19](#page-5-2) and Chile in Table [20](#page-5-4) This would be needed regardless of the compute chosen.

#### Table 13: Implementation with Intel Xeon for full Rubin Observatory

<span id="page-5-5"></span>

#### Table 14: Implementation with Rome for USDF

<span id="page-5-1"></span>

Table 15: Implementation with Intel Xeon for USDF

<span id="page-5-6"></span>

Table 16: Implementation with Intel Xeon for Chile Compute

<span id="page-5-7"></span>

#### Table 17: Implementation with AMD Rome for Chile Compute

<span id="page-5-3"></span>

Table 18: Total storage cost estimate for operations of Rubin Observatory USDF and CHile

<span id="page-5-0"></span>

#### Table 19: Total storage cost estimate for operations at USDF

<span id="page-5-2"></span>

Year (all in M\$)	2024	2025	2026	2027	2028	2029	2030	2031	2032	2033
Fast Storage USDF	\$0.06	\$0.07	\$0.09	\$0.03	\$0.04	\$0.09	\$0.09	\$0.11	\$0.06	\$0.05
Normal Storage USDF	\$3.96	\$3.86	\$4.20	\$4.19	\$4.89	\$8.13	\$8.05	\$8.41	\$8.42	\$8.41
Latent Storage USDF	\$3.18	\$4.69	\$5.37	\$4.72	\$5.26	\$7.90	\$9.40	\$10.09	\$9.44	\$9.44
High Latency Storage USDF	\$0.73	\$1.13	\$1.56	\$2.00	\$2.43	\$2.86	\$3.28	\$3.72	\$4.15	\$4.58
Total (M\$)	\$7.92	\$9.73	\$11.22	\$10.94	\$12.62	\$18.97	\$20.83	\$22.33	\$22.05	\$22.47

Table 20: Total storage cost estimate for operations in Chile **Note** the latent storage here is 0.3 of the USDF latent Storage

<span id="page-5-4"></span>





Table [21](#page-6-0) gives the annual cost of hosting compute in NCSA. This includes purchasing racks to house new nodes etc.

> Table 21: Overheads(NCSA) per year based on number of cores in Table [25](#page-7-0) and costs in Table [30](#page-10-1) assuming Xeon density from Table [29](#page-9-1).

<span id="page-6-0"></span>

### For Chile Table [22](#page-6-1) gives the cost of hosting in Chile (las Serena).

Table 22: Overheads(Chile) per year based on number of cores in Table [25](#page-7-0) and costs in Table [30](#page-10-1) assuming Xeon density from Table [29.](#page-9-1)

<span id="page-6-1"></span>

#### <span id="page-6-2"></span>For Chile the rack costs are outlined in Table [23.](#page-6-2)

Table 23: This table details the cost per rack which is added in Table [22.](#page-6-1)



<span id="page-6-3"></span>For Chile the power costs are outlined in Table [24.](#page-6-3)

Table 24: Cost is estimated to increase 5-10% every 2-3 years





### Various other inputs to ops costing are given in Table [25.](#page-7-0)

Table 25: Various inputs for deriving costs in operations - these drive the costs in Table [5.](#page-2-1) This is based on Table [13](#page-5-5), Table [18](#page-5-0)

<span id="page-7-0"></span>

## **5.1.1 Alternative costing for Chile**

The alternative cost for Chile including Xeons and Qserv but with a similar storage model is given in Table [26](#page-8-2).

Table 26: This table pulls together all the information in a high level summary for Chile operations - in this table Xeon pricing(see Table [16\)](#page-5-7) is used since that is the more expensive but better known option. Price factors, defined in Table [27](#page-8-1) are applied in all cases - other input values come from Table [25](#page-7-0), Table [20](#page-5-4).

<span id="page-8-2"></span>

# <span id="page-8-0"></span>**6 Models**

## **6.1 Sizing model**

An exhaustive and detailed mode is provided in [\[LDM-138;](#page-25-0) [LDM-144\]](#page-25-1) - here we concentrate on the needs for the final years of construction. We explore the compute and storage needed to get us through commissioning and suggest a 2023 purchase for DR1,2 processing which could be pushed to operations.

<span id="page-8-1"></span>Table [27](#page-8-1) gives the annual requirements for the next few years.



Table 27: Various inputs for deriving costs - 2019 represents current holdings.





## **6.2 Compute and storage**

<span id="page-9-0"></span>We which to base our budget on reasonable well know machines for which we have well know prices. Table [29](#page-9-1) gives an outline of a few standard machines we use and a price. This table also gives a FLOP estimate for those machines. Table [28](#page-9-0) gives costs for different types of storage we will require various latency for different tasks and those have varying costs. These tables are used as look ups for the cost models in Section [2](#page-0-0)

Table 28: Storage types and costs used as inputs used for calculations

Storage type	Estimate(SLAC)	Estimate(NCSA)		
fast - NVMe (50GB/seach) / TB	\$400.00	\$1,000.00		
normal - SATA GPFS file systems/ TB	\$133.00	\$135.00		
latency - slower but on disk	\$133.00	\$45.00		
high latency - very slow - on tape	\$15.65	\$25.00		

In Table [28](#page-9-0) we should consider for NVMe for each TB with file system servers two DDN NVMe box with GPFS servers. The price is based on the TOP performer with best price . The Normal price is for each TB with file system disks and servers locally attached to production resources.

In the latency and high latency prices are only at NCSA: for each TB with file systems and all people/services. The complete service not usually attached. S3 bucket type. Can be mounted if needed but not for production worthy speeds. The complete service with data flowing to tape using policies.

<span id="page-9-1"></span>

#### Table 29: Machine types and costs used as inputs used for calculations



<span id="page-10-1"></span>There is also an associated running cost for machines included in the total cost of ownership. These overheads are listed in Table [30](#page-10-1).



#### Table 30: Overhead costs per rack

# <span id="page-10-0"></span>**7 Sizing inputs**

The following simplified sizing was used to give the input sizes for the cost model in Section [2](#page-0-0). The storage sizes are given in Table [33](#page-16-0) and Table [34](#page-17-0) while the compute is given in Table [36](#page-20-0) and Table [37](#page-21-0).

## **7.1 Processing Plan**

This model assumes the following processing:

• Precursor data (HSC RC2 and a similarly-sized DESC DC2 subset) is reprocessed each



month during the Construction period using the Data Release Production (DRP).

- A large precursor reprocessing of HSC PDR2 (or equivalent) is completed twice a year. Products from one of these reprocessings will be released as Data Preview 0. One or more of these processings during Commissioning could be devoted instead to ComCam science data for Data Preview 1 or LSSTCam science data in preparation for Data Preview 2.
- Alert Production (AP) processing happens continuously as LSSTCam science images are obtained. AP hardware is purchased in FY22 to support this, presumably on a limited basis during that year and then to whatever extent possible during the first year of the survey.
- Commissioning processing of LSSTCam science images for Data Preview 2 is assumed to happen towards the end of FY22 as a single execution of the DRP. The hardware for this can be purchased early in that fiscal year.
- Annual DRP execution starts at the beginning of LSST Operations Year 2 with the processing for DR2. The hardware for each year's processing must be purchased and ready for use at the beginning of the year, so it is allocated in the tables to the prior fiscal year, when the images for that processing were taken.
- DR1 processing begins after the first 6 months of the survey; the hardware for this can be part of the DR2 purchase during FY23.

Some storage for raw data needs to be in place at the beginning of the fiscal year, but it can be ramped up over the course of the year. As a simplification it is allocated to the fiscal year in which it will be used.

# **7.2 Storage Model**

<span id="page-11-0"></span>

Table 31: Inputs used to calculate storage needs





## **7.2.1 Overview**

This simplified storage model eliminates many details in the previous storage model [\[LDM-](#page-25-2)[141\]](#page-25-2) that end up being insignificant. There are relatively few data products that require significant amounts of fast SSD or slower disk or tape storage; the others complicate the model without giving much insight. In addition, it is assumed that bandwidth is not a significant constraint, other than the distinction between SSD and spinning disk. With the advent of highly-parallel shared and object storage, having large numbers of spindles solely to achieve high bandwidth for certain operations is not thought to be necessary.

Values are computed for the amount of storage expected to be "on the floor" at the beginning of each fiscal year from FY2020 through FY2023 (which is LSST Operations Year 1). Not included is any storage already present at the end of FY2019 holding past data.

Key scientific and algorithmic assumptions made include:

- All significant intermediates and data products generated by Data Release Production processing need to be kept on filesystem disk until the DRP is complete. Some scratch space is provided to hold small, temporary intermediates. If some intermediates could be removed during DRP when it is known they will no longer be needed, some space savings could be realized.
- HSC RC2 processing is representative of the outputs that DRP will generate. In particular, the presence or absence of "heavy footprints" is assumed to be correct. The coadd storage is doubled to account for an additional "good-seeing" coadd along with the existing "deep" coadd.
- Processed visit images (PVIs) and catalogs in Parquet format start on "normal" filesystem



disk but then move to object storage at the completion of the DRP, with lossy compression of the PVIs at that time. This is in accordance with [RFC-325](https://jira.lsstcorp.org/browse/RFC-325), although the relevant LCR has not yet been approved. Object storage is expected to be cheaper and more scalable for read-only data products; filesystem storage is used for data that is being generated or modified.

- Raw images and coadd images are only temporarily stored on filesystem disk and are then rapidly moved to object storage, where they are retained.
- Intermediates like warped images for coaddition are not survey data products and do not need to be kept beyond the end of the DRP and subsequent QA.

All data is backed up to tape permanently, including annual snapshots of filesystems. Any incremental backups are assumed to be reusable or otherwise purged and hence not significant.

## **7.2.2 Parameters**

The key parameters in Table [31](#page-11-0) are described below.

The numbers of Objects, Sources, and ForcedSources are taken from [LSE-81,](#page-25-3) with the FY2022 numbers reduced by a factor of 2/12 to account for the anticipated 2 months of on-sky science validation time for LSSTCam before the survey begins. These numbers are ultimately based on models for stars in the galaxy and galaxies in the universe that are dependent on the limiting magnitude achieved in each year.

The numbers of science users are estimates, using "Stack Club" users and Commissioning users for FY2020 and 2021, followed by US science users in FY2022 and FY2023 for Data Preview data. The bulk of US science users are not expected to arrive until after Data Release 1 at the beginning of FY2024.

Storage per science user is estimated based on today's usage at NCSA, scaled up as users become more active, and approaching the number given in [LSE-81](#page-25-3) as Operations begins. Note that it is expected that there will be a wide distribution of usage by user, with some using almost none and some using much more than their proportional share.

The LSSTCam image size is uncompressed and includes overscan, 4 bytes of raw data per



pixel, and both science and corner rafts.

The raw image compression factor was measured on simulated LSST images. The lossy image compression factor for processed visit images is the ratio between the lossy-compressed file size (estimated at 1/6 of uncompressed) and the lossless-compressed file size (estimated at 66% of uncompressed). Note that PVIs do not compress losslessly as well as raw images due to their floating point planes.

The number of observing nights per year and the number of visits per night are maximal estimates. 2 images per visit is still the baseline and a possibility that must be accounted for. The number of calibration images per day was derived from the calibration plan.

Two complete all-sky coadds are assumed, one for "good seeing" and one deep.

As stated above, the number of LSSTCam science images is scaled by 2/12 for FY2022 given the length of science validation time. The number of test images, taken on test stands, is estimated as a ramp up to the full science cadence. The numbers of engineering (unprocessed) and calibration images are estimated as ramping-down fractions of the number of science and test images, with calibration images ending at the number per day given previously.

Sizes of rows in various data product tables are taken from [LDM-141,](#page-25-2) which was in turn derived from the DPDD.

Qserv replicates its data for fault tolerance; a typical replication factor is selected here.

### **7.2.3 Data Product Sizing**

Images and the results of processing them are the dominant factor controlling tthe storage sizing which is outlined in Table [32.](#page-16-1) Precursor survey and LSSTCam images are the largest; ComCam, at less than 5% of the size of LSSTCam and with little on-sky science time is negligible, as is LATISS, which is less than 1% of the size of LSSTCam, though it has considerable on-sky time.

The sizing of the Alert Production Database (APDB) is based on experiments in [Salnikov](#page-25-4) [\(DMTN-](#page-25-4)[113\)](#page-25-4) which found that 57,000 visits took 4.5 TB including indexes. A simple linear scaling to a full year's visits was performed, with half that purchased in 2020 for large (but not full) scale



testing.

HyperSuprime-Cam (HSC) RC2 is a relatively small dataset used for monthly processing tests, but it is highly representative of the currently-known DRP work and so is used as the basis for scaling. The size of the input images was taken from [Wood-Vasey et al.](#page-25-5) ([DMTN-091\)](#page-25-5); the size of the outputs (image and Parquet/other non-image files) was measured from the latest execution. A similar size dataset based on DESC DC2 is assumed to be being used for an additional monthly processing test. Note that this is a very small subset of the full DESC DC2, which is expected to cover 300 square degrees to 10-year LSST depth (approximately 1000 epochs per point on the sky). The full DESC DC2 is not currently scheduled to be reprocessed by the construction team. Instead, twice-a-year processings of the full HSC SSP PDR2 dataset (including PDR1) are assumed to occur. The size of this dataset was measured on disk; it is 2,564,358 CCD images, each at 18.2 MB (approximately three times the size of PDR1 alone).

Output sizes are assumed to scale linearly with input size, and by the same factor for each instrument, except for coadds which scale by the sky area processed. While the Object catalog ought to be proportional to sky area as well, its size is expected to be dominated by Source and ForcedSource, so we conservatively make them all proportional to input size (visits) for the precursor data where we do not have object count estimates. For LSSTCam, we use the catalog row estimates to derive Qserv table sizes, but the Parquet file sizes are scaled based on HSC, as they may differ from the Qserv schema.

Scratch space is set at 10% of the output image storage for LSSTCam processing; it is assumed to be already present for precursor processing.

Qserv Czar fast (SSD) storage is assumed to be used for the primary Object table; additional space for the so-called "secondary index" mapping object identifiers to spatial chunks is negligible in comparison.

The main Qserv database storage is based on the Parquet file sizing for precursor data and on the estimated numbers of Objects, Sources, and ForcedSources for LSSTCam data.

Note that no space is explicitly reserved for Qserv query result storage.

An additional 20% disk and tape storage is added to account for all other needs.



<span id="page-16-1"></span>

Table 32: Inputs on dataset sizes used to calculate storage needs

### **7.2.4 Storage Sizing**

Finally, storage is allocated to specific types as shown in Table [33.](#page-16-0) Fast storage (SSD) is used for the APDB and Qserv Czar, which accumulates data from year to year until Data Releases are retired. Normal storage is used for the datasets labeled as such, including output images (initially), output catalogs, and scratch. Local Qserv storage is used for Qserv catalogs. It is assumed that precursor data will be removed from Qserv once LSST data is available, but the LSST data accumulates from year to year. Object storage is used for raw images, lossycompressed output images, lossless-compressed coadd images, and output catalogs. It also accumulates from year to year. Tape is used for long-term archiving of all the raw images, all the data products (both filesystem and object store), and filesystem backups. Again, this accumulates from year to year.

Note that no replication is assumed in the object store.

Table 33: On floor LDF storage estimates based on Table [32](#page-16-1) and Table [31](#page-11-0)

<span id="page-16-0"></span>





An additional table (Table [34](#page-17-0)) gives the storage needs in the Chilean Data Access Center (DAC). This comprises Qserv fast and local storage plus the data products in object storage. Since no DRP computation occurs in Chile, no "normal" filesystem disk is required. Chilean user home directories are assumed to be negligible at this level.



<span id="page-17-0"></span>

# <span id="page-17-1"></span>**7.3 Compute Model**





### **7.3.1 Overview**

This simplified computing model (Table [35\)](#page-17-1) divides computation into three classes: Data Release Production (DRP), Alert Production, and LSST Science Platform (for the US DAC, Chilean DAC, and LSST staff internal use). Calibration Products Production is assumed to be negligible.

The pipelines have advanced considerably in terms of fidelity and science performance since



the previous computing model[[LDM-138](#page-25-0)] was developed. Scaling compute needs based on an execution of the nascent DRP pipeline on HSC PDR1 data and nightly executions of the nascent ap\_pipe pipeline on HiTS2015 data is thus appropriate, but the fact that several steps are still missing from these pipelines must be taken into account.

Elapsed times are measured on existing hardware and converted into core-hours on a nominal CPU (Intel Xeon E5-2680v3 at 2.50 GHz). For example, if a pipeline running on precursor data took an average of one hour on a 32-core nominal CPU, 32 core-hours would be used as its compute requirement. This estimation methodology incorporates all I/O, memory bandwidth, cache miss, and other overheads into the core-hour measurement, simplifying calculations. Note that the nominal CPU does not evolve with time; if future CPUs do more work per core, the actual core-hours may be less than estimated here.

Scaling to other CPUs of the same architecture is based on the ratios of nominal GHz clock rates and core counts. For different architectures (e.g. Rome), the scaling is based on the ratio of industry-reported achievable FLOPS for the two architectures.

Key scientific and algorithmic assumptions are:

- DRP compute time is proportional to the input data size (or, equivalently, the number of visits). While certain tasks are undoubtedly proportional to sky area or number of Objects, overall the pipeline elapsed times are a better fit to the number of visits. Some of this may be because the Object density increases as the number of visits to the same sky patch increases.
- HSC PDR1 processing is generally representative of the final DRP, with an allocation for future additional steps as described below.
- Qserv node counts should remain proportional to the size of data loaded into the database in order to maintain sufficient disk bandwidth and query processing capability, but the proportionality constant changes with time as new generations of system bus with greater bandwidth become available.
- The US DAC LSP is sized at 10% of the DRP compute budget in core-hours, readjusted to be spread over an entire year. The Chilean DAC LSP is sized at 20% of the US DAC (as in [LDM-138\)](#page-25-0). The LSST staff LSP is sized at 10% of the US DAC.



### **7.3.2 Parameters**

The key parameters in Table [35](#page-17-1) are described below.

HSC PDR1 was executed on the NCSA verification cluster, which uses the nominal CPU. The Alert Production executes on Kubernetes nodes, which are a bit slower; to be conservative, this is neglected.

The most recent run of DRP on HSC PDR1 data is described at [https://confluence.lsstcorp.](https://confluence.lsstcorp.org/x/WpBiB) [org/x/WpBiB](https://confluence.lsstcorp.org/x/WpBiB). The input data size is measured; note that the input data files are losslesscompressed. Most jobs (but not most of the time) could run on relatively small-memory machines with 24 cores and 5 GB RAM per core. The largest and longest-running jobs, however, required up to 4 times as much memory, using half or a quarter of the cores. To be conservative, we assume that half the cores were used for the large-memory jobs. The percentage of DRP core-hours that will need to execute on large-memory nodes is estimated.

Since the HSC PDR1 processing did not include several steps from the Science Pipelines Design document [\[LDM-151\]](#page-25-6) such as image differencing and full multi-epoch characterization, the core-hours used are scaled up to the expected pipeline consumption. Note that these algorithmic adjustments are multiplicative.

The SQuaSH system reports the execution time of ap\_pipe in seconds per CCD. A mean was taken over all processed CCDs, and it was assumed that each CCD is processed on a single core. These CCDs are from DECam, which is half the size of an LSST CCD, so the total time is doubled. A factor is added to account for additional steps like differential chromatic refraction compensation and false positive detection that are not well-represented in the current pipeline. Multiplying by the number of LSSTCam science CCDs gives the total number of corehours per visit.

The amount of Qserv data that can be handled by one node is estimated based on the amount of disk that can be scanned in 12 hours at an aggregate rate of 1 GB per second. (Since the Qserv data replicas are not all anticipated to be accessed at the same rate, this is a conservative estimate.)



### **7.3.3 Data Release Production**

The number of nominal core-hours per TB of input data is multiplied by the precursor (HSC RC2 and DESC DC2 subset for 12 months and HSC PDR2 twice a year) and LSSTCam input data sizes (with lossless compression) to determine the total number of core-hours needed in each year. This is shown in Table [36.](#page-20-0) Approximately one-third of these core-hours need to be provided by small-memory (4-5 GB/core) machines; the other two-thirds need to come from large-memory (8-20 GB/core) machines.

<span id="page-20-0"></span>



### **7.3.4 Alert Production**

The core-hours per visit are divided by the minimum visit length (30 sec plus 1 sec shutter motion plus 2 sec readout) to give the minimum number of cores needed to keep up with image taking. This is shown in Table [36](#page-20-0). These cores are expected to be provided over multiple "strings" of nodes. Note that the current AP design is not readily able to take advantage of more than one core per CCD.

#### **7.3.5 LSST Science Platform**

LSST Science Platform needs for US DAC science users are derived as 10% of the DRP corehour requirement and are shown in Table [37.](#page-21-0) The LSP core-hours are assumed to be spread over a year, giving the total number of nominal cores needed in the DAC. Peak loads are expected to be handled by "borrowing" elastically from the DRP compute pool.

As a reasonableness check, the number of cores per science user is computed, but it must be noted that an oversubscription factor needs to be taken into account since not all users are expected to be simultaneously active.

Similar computations for the Chilean DAC (at 20% of the US DAC) and the LSST staff LSP (at



10% of the US DAC) are also in Table [37.](#page-21-0)

The number of Qserv nodes needed is computed from the storage devoted to it and the storage per node number. Note that staff use of Qserv is taken into account by loading the Data Release products into an internal-only Qserv instance and then making that instance part of the DAC at Data Release, so the compute sizing is part of the US DAC.

<span id="page-21-0"></span>

<b>US DAC</b>	units	<b>FY20</b>	<b>FY21</b>	<b>FY22</b>	<b>FY23</b>	<b>Notes</b>
LSP cores	cores			128	517	10% of DRP, over a year
Oserv nodes	nodes			14	95	
LSP cores/ science user	cores/user			0.03	0.10	includes oversubscription
Chilean DAC	units	FY20	FY21	<b>FY22</b>	<b>FY23</b>	Notes
LSP cores	cores			26	103	20% of US DAC
Oserv nodes	nodes			14	95	
Staff LSP	units	<b>FY20</b>	FY21	<b>FY22</b>	FY23	Notes
LSP cores	cores			13	52	10% of US DAC

Table 37: Compute needs for the Science Platform instances

### **7.3.6 DES Comparison**

As another check on the model, core-hour figures for Dark Energy Survey (DES) processing were obtained. These are given in Table [38](#page-21-1). The CPUs used for single-frame and coadd processing had slightly slower clock rates but better bandwidths and expected instructions per clock performance, so they were considered equivalent to our nominal core. The CPUs used for Multi-Object Multi-Band Fitting and Single-Object Fitting (MOF/SOF) included a large contribution from the Blue Waters machine at NCSA. Those CPUs (AMD 6276) are somewhat older and were estimated at 0.245 nominal cores.

The single-frame processing measured number of 5.2 core-hours per visit compares well with the 5.4 core-hour per visit parameter used in our sizing model. Similarly, the overall DES compute figure of 21,000 core-hours per terabyte is virtually identical to our estimate (including the factors for additional steps).

<span id="page-21-1"></span>

#### Table 38: Comparison with DES compute



## **7.4 Operations Sizing**

Five tables use some of the parameters from the above model to project LSST storage and compute needs throughout the 10 years of Operations.

## **7.4.1 Storage in Operations**

The Object, Source, and ForcedSource numbers in Table [39](#page-23-0) are taken from [LSE-81](#page-25-3), as before. The number of science users and storage per user is ramped up. Note that the number of images needing storage and processing grows linearly with time. Table row sizes are taken from [LDM-141](#page-25-2); they include growth over time as columns are added.

The dataset sizes in Table [40](#page-23-0) are calculated using the same formulas and proportionality constants as in Table [32](#page-16-1).

The on-the-floor storage estimates in Table [41](#page-24-0) include fast (SSD) storage for the APDB and Qserv Czar, with the latter being sized for three Data Releases (two being served and one being prepared).

"Normal" filesystem storage holds raw images, data products, scratch space, Qserv data prior to loading, science user workspace, and a 20% allocation for everything else.

Qserv local storage holds catalogs for three Data Releases.

Raw images (lossless-compressed) are written immediately to object storage, as are Parquetformat catalogs. PVIs are lossy-compressed and placed in object storage. The complete set of raw images is available, whereas the catalogs from only the last two Data Releases and the one in preparation are kept, and the PVIs from only the last Data Release and the one in preparation are online.

All data products and new raw images for each Data Release are copied to tape, but scratch space and the Qserv-schema catalogs are not.

Table [42](#page-24-0) extracts the Qserv and object store sizing needed to populate the Chilean DAC with a copy of the data products and raw images.

Table 39: Inputs used to calculate storage needs during Operations



#### Table 40: Dataset sizes used to calculate storage needs during Operations

<span id="page-23-0"></span>

#### Table 41: On floor LDF storage estimates during Operations



#### Table 42: On floor Chile storage estimates during Operations



#### Table 43: Compute needs during Operations

<span id="page-24-0"></span>



### **7.4.2 Compute in Operations**

The DRP compute sizing in Table [43](#page-24-0) follows directly from the size of the input data to be processed. The number of cores for Alert Production does not change with time. The DAC and staff LSP instances are sized based on the assumed percentages of DRP compute. The amount of Qserv data that can be handled by a node is assumed to grow with time, doubling every four years (PCI Express has gone from 1.0 GB/sec to 16 GB/sec between 2003 and 2019). The number of Qserv nodes is calculated by dividing each Data Release's storage by the storageper-node figure for its year; older nodes are assumed to be retired.

# **A References**

# **References**

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- <span id="page-25-6"></span>**[LDM-151]**, Swinbank, J.D., et al., 2017, *Data Management Science Pipelines Design*, LDM-151, URL <https://ls.st/LDM-151>
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# **B Acronyms**





